## A Strategic Analysis of Information Sharing Among Cyber Attackers

Anindya Ghose Leonard N. Stern School of Business New York University 44 West 4th Street, Suite 8-94 New York, NY 10012-1126, USA E-mail: aghose@stern.nyu.edu Tel.: +1 (212) 998-0807, Fax: +1 (212) 995-4228

Kjell Hausken Faculty of Social Sciences University of Stavanger N-4036 Stavanger, Norway E-mail: kjell.hausken@uis.no Tel.: +47 51 831632, Fax: +47 51 831550

Version: August 24, 2006

Keywords: Cyber war, hacking, attack, defense, conflict, contest success function, security investment, information sharing, security breaches.

## Acknowledgements

We thank Ross Anderson for suggestions that led to some additional analysis, and Dean Abrams for inquiring about this phenomenon which led to the completion of this paper.

## Abstract

One firm invests in security to defend against cyber attacks by two hackers. Each hacker chooses an optimal attack, and they share information with each other about the firm's vulnerabilities. Each hacker prefers to receive information, but delivering gives competitive advantage to the other hacker. We find that each hacker's attack and information sharing are strategic complements while one hacker's attack and the other hacker's information sharing are strategic substitutes. The attack is inverse U-shaped in the firm's unit defense cost, and reaches zero, while the firm's defense and profit decrease, and the hackers' information sharing and profit increase. The firm's profit increases in the hackers' unit cost of attack, while the hackers' information sharing and profit decrease. Our analysis also reveals the interesting result that the cumulative attack level of the hackers is not affected by the effectiveness of information sharing between them and moreover, is also unaffected by the intensity of joint information sharing. We also find that as the effectiveness of information sharing between hackers increases relative to the investment in attack, the firm's investment in cyber security defense and profit are constant, the hackers' investments in attacks decrease, and information sharing levels and hacker profits increase. In contrast, as the intensity of joint information sharing increases, while the firm's investment in cyber security defense and profit remain constant, the hackers' investments in attacks increase, and the hackers' information sharing levels and profits decrease. Increasing the firm's asset causes all the variables to increase linearly, except information sharing which is constant. We extend our analysis to endogenize the firm's asset and this analysis largely confirms the preceding analysis with a fixed asset.

## 1. Introduction

The increasing ubiquity of the Internet provides cyber hackers more opportunities to misappropriate or corrupt an organization's data resources. There are many well known examples of cyber-hacking such as Egghead.com which faced a massive backlash from its customers after being hacked in 2000 and Travelocity where hackers exposed the personal information of thousands of customers who had participated in an online promotion in 2001. Even well established firms like Citibank, Microsoft and NASA, among others have been targeted and hacked by cyber perpetrators.

For a while now, it has been recognized that a key factor required to improve computer and information security is the gathering, analysis and sharing of information related to successful, as well as unsuccessful attempts at, computer security breaches. To encourage information sharing among organizations, the US federal government has encouraged the establishment of Security Based Information Sharing Organizations (SB/ISOs) of various kinds, such as Information Sharing & Analysis Centers (ISACs), CERT, INFRAGARD, etc. Similar initiatives have been taken in other parts of the world.

There are several positive aspects to sharing information about security incidents and vulnerabilities. These benefits include both the prevention of further security breaches in the future (e.g., identifying and repairing vulnerabilities in their information security systems) as well as increased sales resulting from more effective security products and better security reputation among consumers. Questions on information sharing, economic incentives and social welfare similar to those noted above have been previously studied in the context of other organizations. Of particular relevance is the extensive literature on trade associations (TAs). Previous relevant work includes in the literature on oligopolies, cooperative relationships, joint ventures, and trade associations (Gal-Or 1986, Kirby 1988, Novshek and Sonnenschein 1982, Shapiro 1986, Vives 1990). More recently, information sharing among firms to defend against cyber attacks has been analyzed by Gordon, Loeb and Lucyshyn (2003), Gal-Or and Ghose (2005), and Hausken (2006). The focus of Gordon et al. (2003) is on how information sharing affects the overall level of information security. They highlight the tradeoff that firms face between improved information security and the potential for free riding, which can lead to under-investment in security expenditures. While Gordon et al. (2003) focus on the cost side

effects of security breaches and information sharing, Gal-or and Ghose (2005) focus on the demand side effects and highlight the strategic implication of competition in the product market on information sharing and security technology investment levels.

Information sharing among hackers is different. Firms subject to attacks would naturally prefer that such information sharing and cooperation among hackers does not take place. As part of their activities, hackers compile information about firms' vulnerabilities and defense strategies, attempt to gain access to the information the firms collect about their security breaches and share the information among each other.

What motivates hackers to hack? Howard (1997) identified the five possible objectives for hackers- financial gain, a desire for challenges, political gain, a desire to cause destruction, and leisure activities.<sup>1</sup> Some research has pointed out that greed, power, and revenge are superseding curiosity and other benign motivations (Jordan and Taylor 1998). Similarly some work in the sociology and computer science literature (Raymond 2001) on motivation of hackers has pointed out that because of issues such as reputation and competition amongst hackers in order to get more recognition within the community, hackers may have incentives to not share information with each other. Indeed hackers actively compete with one another to write the best software, frequently one-upping each other in displays of coding prowess. Combining hacker egos with their practice of sharing fixes amongst one another and the sheer joy hackers take from hacking, one can clearly understand how reputation plays a large role as pointed out by Ritchie (2000). On the other hand, there has also been discussion about how hackers generally don't believe in keeping secrets and are quite keen to share information (Brunker 1998)<sup>3</sup> and not necessarily compete for reputation from a sociological view (Risan 2000). In sum, we believe there are reasons for both viewpoints but the key point is that hackers do share information with each other.

This paper analyzes two hackers who may share information about a firm's vulnerabilities, in addition to designing the size of their attacks. The firm invests in information technology

<sup>&</sup>lt;sup>1</sup> He further developed a taxonomy of the hacking process involving the attacker, tools, access, results and objectives of attacks. Kjaerland (2005) extended to account for the target or victim of an attack.

<sup>&</sup>lt;sup>3</sup> Hackers: Knights-errant or knaves? http://msnbc.msn.com/id/3078783/

security to defend against the attacks.<sup>4</sup> Naturally, each hacker prefers to receive information from the other hacker, but is reluctant to deliver information, though there are benefits from joint information sharing. We assume that both hackers and the defending firm are strategic actors. The opponent does not have a given, fixed, or immutable strategy, which has been common in much of prior research in information security. The absence of an assumption about a fixed threat, or a fixed defense, enables a much richer analysis.

In a related stream of work, Png, Tang and Wang (2005) focus on the strategic interaction among end-users and between users and hackers in a setting with a continuum of user types, and show that users' effort in fixing depends on hackers' targeting and vice-versa. Prior work (Choi, Fershtman, and Gandal 2004, Nizovtsev and Thursby 2005; Arora, Caulkins, and Telang 2005) has examined issues such as incentives for security specialists to disclose security flaws and provide the appropriate patches Other work (Cavusoglu, Mishra and Raghunathan 2005), and Anderson (2001)) has examined issues such as users' incentives to invest in intrusion detection systems observed that many systems fail not for technical reasons so much as from misplaced incentives.

While the two hackers may be interpreted as two individual agents seeking to exploit a firm, one can also interpret them as two firms who decide to gang up on a third rival firm to exploit its asset. The reason for this broad interpretation is that the terms "attack" and "defense" can be understood as metaphors. As Hirshleifer (1995) puts it, "falling also into the category of interference struggles are political campaigns, rent-seeking maneuvers for licenses and monopoly privileges (Tullock 1967), commercial efforts to raise rivals' costs (Salop and Scheffman 1983), strikes and lockouts, and litigation – all being conflictual activities that need not involve actual violence." Attack and defense are subcategories of attack-oriented and defensive competition. In this paper we use the more precise terms such as attack and defense, which can be substituted with synonyms such as struggle, conflict, battle, etc.

Information sharing is technically complicated to analyze. Since our objective is to analyze information sharing among hackers, and not how the firm defends in a possibly different manner against two hackers who may be configured differently, we confine attention to two equivalent

<sup>&</sup>lt;sup>4</sup> We model competition between hackers and between hackers and the firm, but not between firms. The firm can be thought of as being a federal government institution or a public sector agency and hence competition at

hackers which makes the analysis tractable. This allows analyzing more specifically, the tradeoffs between information sharing and cyber attack, which are the two strategic choice variables for each hacker, in interaction with a defending firm which has one strategic choice variable-the level of investment in its security.

One main difference between cyber attacks and information sharing is that the former requires costly funding, planning, sustained effort through time, involving buildup of infrastructure, culture, and competence, while the latter may be more or less costless for each hacker aside from the time spent in transferring the information. Attackers in the cyber era of course have information storage capacity, and they may even have compiled and stored information in an organized and secure manner. Deciding to share information with another hacker may not involve more than sending an electronic mail, or storing the information on a disk and delivering it. In other words, designing cyber attacks entail explicit costs on part of the hackers since all investments are generally costly, while information sharing does not have explicit costs, although the competitive advantage given to a hacker can be construed as an implicit cost.

For firms defending against cyber attacks Gal-Or and Ghose (2005) find that security investments and information sharing are strategic complements. Gordon et al. (2003) find that when firms share information, each firm has reduced incentives to invest in information security. Hausken (2006) finds that security investments and information sharing are strategic substitutes. This paper examines these issues for two hackers. An important difference is that whereas two firms incur information leakage costs when sharing information about their vulnerabilities and security breaches, two hackers sharing information about a firm's vulnerabilities do not incur costs in the same manner. In fact, as a hacker shares information about a firm with another hacker, the firm itself is the main entity suffering through incurring information leakage costs which are possible and often likely in all information transfers.

The work by Gordon et al. (2003), Gal-Or and Ghose (2003, 2005), Hausken (2006), and also this paper, assume that information can be scaled along one dimension. Gordon et al. (2003) refer to a portion, which is a number between zero and one, of a firm's computer security information that it may decide to share with the other firm. Similarly, Gal-Or and Ghose (2005) normalize the amount of security information being shared so that it always lies between 0 and

the firm level may not be critical in such situations.

1. Generally, information is multi-faceted, of different kinds, and with different degrees of importance for different purposes. In our context, a one-dimensional concept of information can be interpreted to mean that different kinds of information are given different weights according to their relative importance.

Section 2 presents the model. We start with the case when the firm's asset is exogenous. Section 3 analyzes the model and presents some numeric analysis to provide insights. In section 4 we endogenize the firm's asset, so that the firm makes a tradeoff between producing the asset and security investment needed for defending it. Section 5 assumes an upper constraint on each hacker's fraction of the asset before it is closed down and provides some basic insights. Section 6 discusses various extensions and limitations. Section 7 concludes with some discussion of the implications of our study.

## 2. The model

Consider one firm with an asset *r* and two equivalent hackers *i* and *j* launching cyber attacks to acquire portions of the firm's asset. The firm invests *t* to defend its asset, and the defense expenditure is *f*, where  $\partial f / \partial t > 0$ . The hackers invest  $T_i$  and  $T_j$ , respectively, to attack the asset, with attack expenditures  $F_i$  and  $F_j$ , respectively, where  $\partial F_i / \partial T_i > 0$  and  $\partial F_j / \partial T_j > 0$ . For simplicity, we consider a linear function given by f=ct,  $F_i=\theta T_i$ ,  $F_j=\theta T_j$ , where 1/c is the efficiency of cyber defense and  $1/\theta$  is the efficiency of cyber attack. This means that *c* and  $\theta$  are the investment inefficiencies. They can also be interpreted as unit transformation costs. An attack means attempting to break through the security defense of the firm in order to appropriate something of value to the firm<sup>5</sup>. For simplicity, we assume risk neutral actors, which does not change the nature of the argument. The expenditures *ct*,  $\theta T_i$ ,  $\theta T_j$  can be interpreted as expenses in capital and/or labor.

The cyber contest between the firm and the two hackers for the asset r takes the common ratio form (Tullock 1980). In the absence of information sharing, we consider the following contest success function:

<sup>&</sup>lt;sup>5</sup> This could be customer related information, business strategy information or accounting related information.

$$g = \frac{t}{t + T_i + T_j},$$

$$G^i = \frac{T_i}{t + T_i + T_j},$$

$$G^j = \frac{T_j}{t + T_i + T_j},$$
(1)

where g,  $G^i$  and  $G^j$  are the fractions of the firm's asset that the firm, hacker *i*, and hacker *j* retain. The three fractions sum to one,  $g+G^i+G^j=1$ . As expected, the first fraction satisfies  $\partial g / \partial t > 0$ ,  $\partial g / \partial T_i < 0$ ,  $\partial g / \partial T_j < 0$ . That is, the firm benefits from its own security investment, and suffers from the attacks launched by the two hackers. Similarly, the second fraction satisfies  $\partial G^i / \partial t < 0$ ,  $\partial G^i / \partial T_i > 0$ ,  $\partial G^i / \partial T_j < 0$ , and similarly for the third fraction. These fractions imply that each hacker benefits from its own investment, but suffers a loss from the investment made by the firm and an attack by the other hacker on the firm.

Hausken (2006) analyzes two interdependent firms subject to cyber attacks, and shows how increasing interdependence causes increased information sharing and decreased security investment. An attack against one firm gets partly propelled further to the other firm due to the interdependence, and a firm's defense operates both against direct attacks and indirect attacks, where the latter are those proceeding through the other firm due to the interdependence.

With two hackers, the interdependence between these is not relevant in the same sense for the following reasons. First, an attack launched by hacker *i* is directed against the firm, and it does not make sense to consider this attack as somehow also proceeding through hacker *j* to the firm. Secondly, the defense set up by the firm operates against the total sizes  $T_i$  and  $T_j$  of the two attacks, where the interdependence between the hackers plays no role.

As a hacker's attack level increases, it compiles more information about the firm's vulnerabilities. Characteristics of the information are the type of firewalls, encryption techniques, access control mechanisms, intrusion detection systems, etc. employed by the firm, the training and procedures of the firm's security experts, the nature of the defense, and the properties of the vulnerabilities. A lot of this information is publicly available.

Assume that hacker *i* decides to share an amount  $S_i$  of information with hacker *j*, and that hacker *j* similarly decides to share an amount  $S_j$  of information with hacker *i*. Naturally, hacker *i* prefers to receive information from hacker *j* since it gets thereby more insight into the firm's vulnerabilities and defensive strengths, which makes it better equipped to launch a successful attack. Analogously, hacker *j* prefers to receive information from hacker *i*. Based on this reasoning, we assume that hacker *i*'s overall attack has a strength equal to  $T_i + \gamma S_j$ , and that hacker *j*'s overall attack has a similar strength of size  $T_j + \gamma S_j$ , where the parameter  $\gamma$  measures the effectiveness of information sharing relative to the attack in the composition of the overall attack.  $\gamma$  varies between 0 and 1. With such information sharing, the contest success functions in (1) become equal to

$$h = \frac{t}{t + T_i + \gamma S_j + T_j + \gamma S_i},$$

$$H^i = \frac{T_i + \gamma S_j}{t + T_i + \gamma S_j + T_j + \gamma S_i},$$

$$H^j = \frac{T_j + \gamma S_i}{t + T_i + \gamma S_j + T_j + \gamma S_i}.$$
(2)

These three fractions capture three important consequences from information sharing: First, it incorporates the fact that the firm suffers a loss from information sharing among the hackers. Second, hacker *i* benefits from receiving information from hacker *j*, but suffers a loss in utility from delivering information since hacker *j* gets a competitive advantage. This follows because  $\gamma S_j$  is present in the numerator in the second fraction in (2), while both  $\gamma S_j$  and  $\gamma S_i$  are present in the denominator. Consequently, the hackers can be expected to free ride on each other's information sharing, and to the extent (2) is descriptive, no information sharing will occur. For the second fraction in (2),  $\partial H^i / \partial S_i < 0$  and  $\partial H^i / \partial S_j > 0$ , and analogously for the third fraction  $\partial H^j / \partial S_i < 0$  and  $\partial H^j / \partial S_i > 0$ .

As the hackers share information with each other, synergies emerge. For instance, they discuss the available information, transformation occurs, missing pieces are filled in, and reasoning based on the joint information generates new knowledge. Joint information sharing by the two hackers can thus be expected to generate even deeper insight into the firm's vulnerabilities and defense. To model this effect, we introduce the multiplicative term  $S_iS_j$ , and a parameter  $\lambda$  which scales the intensity of joint information sharing. Multiplicative terms are common to express joint interaction. For example, in epidemiology, the rate by which the number of susceptible individuals (those that are not infected, but are capable of catching the disease and get infected) decreases is proportional to the product of the number of susceptible individuals and the number of infected individuals (Kermack and McKendrick 1927). Similarly, in Lanchester (1916) guerilla warfare the loss rate of each group is proportional to the product of the sizes of the two groups. This product expresses the number of contacts between the two groups.

For a fixed number of individuals in the two groups together, the number of contacts is largest if the groups are equally large, and decreases as the groups grow unequal in size. Similarly, for information sharing, the joint benefit is large when both hackers contribute similarly large amounts of information, and decreases as the contributions grow unequal given that the sum of the contributions remains the same. If one hacker shares substantially, while the other hacker shares very little, the joint benefit is limited since the small amount of information can throw light on, supplement, and complement the large amount of information to a small extent.<sup>6</sup> With such joint information sharing, the contest success functions in (2) can be written as

$$k = \frac{t}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j},$$

$$K^i = \frac{T_i + \gamma S_j + (\lambda/2) S_i S_j}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j},$$

$$K^j = \frac{T_j + \gamma S_i + (\lambda/2) S_i S_j}{t + T_i + \gamma S_i + T_j + \gamma S_i + \lambda S_i S_j}.$$
(3)

We let the two hackers share the benefit of their joint sharing so that each of them receives  $(\lambda/2)S_iS_j$ . Equation (2) expresses the tradeoff each hacker makes between how much information to share and how much to withhold. Intuitively, it prefers to share sufficient information so that it reaps the benefits of joint information sharing. However, if it shares too much information, the other hacker is given an undue advantage because it can get a larger portion of the firm's asset. This places an upper limit on how much information is optimal for a hacker to share.

Gal-Or and Ghose (2005) and Hausken (2006) assume that two firms sharing information about their security breaches incur information leakage costs. Such costs are to a firm's disadvantage.

<sup>&</sup>lt;sup>6</sup> One illustrative example is that  $5 \times 5=25$  is larger than  $2 \times 8=16$ , although 5+5=2+8=10.

Basically, leakage of security vulnerabilities could increase consumers' apprehension of transacting with a firm. As a result some consumers may find it optimal to switch from one firm to its rival, thereby reducing its demand. When two hackers share information they possess about the firm's vulnerabilities, security breaches, and defense characteristics, there can also be some information leakage. However, the costs from such leakage are borne by the firm. Although two hackers' sharing may cause leakage through other networks and channels than when two firms share information, the information about security vulnerabilities pertains to the firm or firms which incur(s) substantial costs when such leakage occurs. For example, the information leak may progress to a firm's rival, which may enable the rival to use this information strategically against its competitor. In contrast it's plausible that such information sharing does not harm the hackers per se. After all, they have much less to lose. Hence, when one hacker shares information with another hacker and a leak occurs, we assume that the hackers incur no leakage costs; these costs affect the firm's payoff.

In accordance with Gal-Or and Ghose (2005), we assume that the leakage costs caused by information sharing by hacker *i* is given by the function  $g^i = \phi_1 S_i^2 - \phi_2 S_j^2 - \phi_3 S_i S_j$ , and analogously leakage costs caused by information sharing by hacker *j*, is given by the function  $g^j = \phi_1 S_j^2 - \phi_2 S_i^2 - \phi_3 S_i S_j$ . Note that these costs sum to  $(\phi_1 - \phi_2)(S_i^2 + S_j^2) - 2\phi_3 S_i S_j$ , where  $\phi_1 \ge \phi_2 + \phi_3$ .

The profits u,  $U_i$ ,  $U_j$ , of the firm, hacker *i*, and hacker *j*, respectively, are given by

$$u = \frac{t}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j} r - ct - \left[ (\phi_1 - \phi_2)(S_i^2 + S_j^2) - 2\phi_3 S_i S_j \right],$$

$$U_i = \frac{T_i + \gamma S_j + (\lambda/2)S_i S_j}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j} r - \theta T_i,$$

$$U_j = \frac{T_j + \gamma S_i + (\lambda/2)S_i S_j}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j} r - \theta T_j.$$
(4)

#### 3. Analysis

The firm's free choice variable is the security defense expenditure level *t*. Hacker *i*'s free choice variables are the level of security attack  $T_i$  and the level of information sharing  $S_i$ . Analogously, hacker *j*'s free choice variables are  $T_i$  and  $S_i$ . The game proceeds as follows. The firm and the

two hackers choose their free choice variables simultaneously and independently to maximize profits. Calculating  $\partial U_i / \partial S_i = 0$  and solving with respect to  $S_i$  gives

$$S_{j} = \frac{t - T_{i} + T_{j} - 2\gamma^{2} + \sqrt{(t - T_{i} + T_{j} - 2\gamma^{2})^{2} - 8T_{i}\gamma^{2}}}{2\gamma}$$
(5)

This leads to the following observation:

**Observation 1.** *(i)* Each hacker's investment in an attack and the level of information sharing are strategic complements such that increasing one causes an increase in the other, and vice versa. (ii) One hacker's attack and the other hacker's information sharing are strategic substitutes such that increasing one causes a decrease in the other, and vice versa.<sup>7</sup>

Note that an increase (decrease) in  $T_j$  on the right hand side of (5) causes an increase (decrease) in  $S_j$  on the left hand side. Similarly, an increase (decrease) in  $T_i$  on the right hand side of (5) causes a decrease (increase) in  $S_j$  on the left hand side.

Observation 1 states that if one hacker were to increase its attack, then it also increases its level of information sharing. However, if a hacker receives more information from the other hacker, then it cuts back on its own attack level. It is useful at this stage to compare this result with prior work that looks at information sharing between two firms. When two firms share information and defend against one external hacker, Hausken (2006) finds that information sharing and security investments of the two firms are strategic substitutes. Gordon et al. (2003) find strategic substitutability under certain assumptions in the sense that when firms share information, each firm has reduced incentives to invest in information security.<sup>8</sup> Gal-Or and Ghose (2003, 2005) find strategic complementarity so that that increased security investment by one firm leads to increased security investment and increased information sharing by its competitor.

<sup>&</sup>lt;sup>7</sup> Ideally one would like to generalize the above observation into a broader result using the implicit function approach after total differentiation of the three first-order conditions and sign the Hessian Matrix using Cramer's Rule. We tried doing that. However, the resultant Hessian Matrix is impossible to sign without making some strong assumption about parameter values. Hence, we do not state this result as a formal Proposition. However, it is important to point out that the qualitative nature of the intuition and insights are the same. Hence, we focus on discussing the results in the form of an Observation and contrast them with results from prior work.

<sup>&</sup>lt;sup>8</sup> Information sharing and security investment are not pure strategic substitutes in Gordon et al.'s (2003) study. They provide necessary and sufficient conditions for information sharing to lead to increased or decreased security investment.

Next we derive the first order conditions. By setting  $\partial u / \partial t = 0$ ,  $\partial U_i / \partial T_i = 0$ ,  $\partial U_i / \partial S_i = 0$ , and thereafter substituting  $T = T_i = T_j$  and  $S = S_i = S_j$  in equilibrium, gives the three first order conditions.

FOC 
$$t: \frac{r(2T+S(2\gamma+\lambda S))}{(t+2(T+\gamma S)+\lambda S^2)^2} - c = 0,$$
 (6)

FOC 
$$T: \frac{2r(t+T+\gamma S)+r\lambda S^2}{2(t+2(T+\gamma S)+\lambda S^2)^2} - \theta = 0,$$
 (7)

FOC 
$$S: 2\gamma(T+\gamma S) + \lambda S(\gamma S - t) = 0.$$
 (8)

Solving with respect to *t*, *T*, *S*, setting  $K=K^i=K^j$  and  $U=U_i=U_j$  in (3) and (4), and inserting into the profit functions *u* and *U* and gives the following interior solution:

$$t = \frac{2r(2\theta - c)}{(2\theta + c)^{2}},$$

$$T = 2c \left(\frac{r}{(2\theta + c)^{2}} - \frac{2\gamma^{2}\theta}{(2\theta - c)^{2}\lambda}\right),$$

$$S = \frac{2c\gamma}{(2\theta - c)\lambda},$$

$$k = 1 - \frac{2c}{2\theta + c},$$

$$K = \frac{c}{2\theta + c},$$

$$u = \frac{r(2\theta - c)^{2}}{(2\theta + c)^{2}} - \frac{8c^{2}\gamma^{2}(\phi_{1} - \phi_{2} - \phi_{3})}{((2\theta - c)\lambda)^{2}},$$

$$U = c \left(\frac{cr}{(2\theta + c)^{2}} + \frac{4\gamma^{2}\theta^{2}}{(2\theta - c)^{2}\lambda}\right)$$
where  $2\theta > c$ ,  $r > \frac{2\gamma^{2}\theta(2\theta + c)^{2}}{(2\theta - c)^{2}\lambda}.$ 
(9)

The hackers' overall attack level, which can also be referred to as the cumulative attack level, as seen from the two numerators in the last two expressions of equation (4), dependent on investment, information sharing, and joint information sharing and equals the following:

$$T_A = T + \gamma S + (\lambda/2)S^2 = \frac{2cr}{(2\theta + c)^2}$$
(10)

13

In equilibrium,  $T_A$  is independent of  $\gamma$  and  $\lambda$ . This also explains why the firm's defense level *t* is independent of  $\gamma$  and  $\lambda$ , and *u* is independent of  $\gamma$  and  $\lambda$  when there are no leakage costs, i.e., when  $\phi_1 = \phi_2 + \phi_3$ . Further, note that the fractions *k* and *K* of the asset, where *k*+2*K*=1, accruing to the firm and to each hacker, respectively, are independent of  $\gamma$  and  $\lambda$ , and of the asset *r*. This leads to the following result.

**Proposition 1** (*i*) The fraction of the asset accruing to each hacker is equal to the ratio of the firm's investment inefficiency, c, to the sum of the investment inefficiencies of all the three actors,  $2\theta+c$ .

(ii) The cumulative attack level  $T_A$  of the hackers is independent of both the effectiveness of information sharing,  $\gamma$ , and the intensity of the joint information sharing,  $\lambda$ .

(iii) Further, the cumulative attack level (i) increases in the inefficiency of the firm's cyber defense at a decreasing rate, (ii) decreases in the inefficiency of the hacker's cyber attack at an increasing rate and (iii) increases in the value of the firm's asset. That is,  $\partial T_A / \partial c > 0$ ,  $\partial^2 T_A / \partial c^2 < 0$ ,  $\partial T_A / \partial \theta < 0$ ,  $\partial^2 T_A / \partial \theta^2 > 0$ ,  $\partial T_A / \partial r > 0$ ,  $\partial^2 T_A / \partial r^2 = 0$ .

An interesting result is that the cumulative attack level  $T_A$  of the hackers is not affected by the effectiveness  $\gamma$  of information sharing between hackers and moreover, is also unaffected by intensity  $\lambda$  of joint information sharing. This result can have some ramifications. For instance, organizations are always worried about the extent to which information sharing between hackers can adversely affect them. Our analysis reveals that the strength of the total attack is only related to the extent of the firm's inefficiency in cyber defense, the hackers' inefficiencies in attack and the asset value of the firm. If it turns out that information sharing between hackers does not enhance the attack intensities, then there is less cause for concern from such information sharing than it is thought to be.

**Proposition 2**. As the extent of inefficiency of the cyber defense of the firm increases towards the extent of inefficiencies of the cyber attack of the two hackers, (that is, as c increases towards  $2\theta$ ), the firm's security investment and the hackers' information sharing decrease toward zero

but remain strictly positive, t>0 and S>0t. Conversely, the hackers' attack level T reduces to zero for a value of c that is strictly below  $2\theta$ .

Observe in (9) how each hacker's inefficiency  $\theta$  is always preceded by 2, while the firm's inefficiency *c* is preceded by 1, that is, it is equal to *c*. This occurs because there are two hackers, but only one firm. Furthermore, the firm does not defend itself when  $2\theta < c$ . When the firm's inefficiency is larger than the double inefficiency of each hacker, then the firm gives up its asset due to its intrinsic inefficiency in protecting its asset. The inequality can also be phrased as  $1/c < 1/2\theta$  or as  $2/c < 1/\theta$ . This means that the defense efficiency 1/c of the firm must be at least 50% of the attack efficiency of each hacker in order for the firm to find it worthwhile to invest in some level of cyber defense. When the firm does not invest at all in defense, the two hackers share the firm's asset and do not share information.

When T = 0 according to Proposition 1, solving the first order conditions for *t* and *S* in (6) and (8) gives a third order equation set which is voluminous to set up and tedious to analyze generally. We thus confine our general discussion to the interior solution. The following proposition specifies the dependence of the six variables *t*, *T*, *S*, *u*, *U*, and *K* on the five parameters *c*,  $\theta$ ,  $\gamma$ ,  $\lambda$ , *r*.

## **Proposition 3.** Assume $\phi_1 = \phi_2 + \phi_3$ . Then,

(i) the firm's investment in cyber defense decreases at an increasing rate in the level of its own inefficiency. Further, the hackers' level of information sharing increases at an increasing rate in the level of the firm's inefficiency in cyber defense. Finally, the firm's profit decreases at an increasing rate in its own inefficiency. That is,  $\partial t / \partial c < 0$ ,  $\partial^2 t / \partial c^2 > 0$ ,  $\partial S / \partial c > 0$ ,  $\partial^2 S / \partial c^2 > 0$ ,  $\partial u / \partial c < 0$ ,  $\partial^2 u / \partial c^2 > 0$ .

(ii) the hackers' level of information sharing and pay-offs decrease at an increasing rate in the level of their inefficiency in cyber attacks. That is,  $\partial S / \partial \theta < 0$ ,  $\partial^2 S / \partial \theta^2 > 0$ ,  $\partial U / \partial \theta < 0$ ,  $\partial^2 U / \partial \theta^2 > 0$ . Further, when  $c/2 < \theta < 3c/2$ , the firm's investment in cyber defense increases at a decreasing rate in the level of the hackers' inefficiency in cyber attacks. That is,  $\partial t / \partial \theta > 0$ ,  $\partial^2 t / \partial \theta^2 < 0$  when  $c/2 < \theta < 3c/2$ .

(iii) the firm's investment in cyber security is independent of the effectiveness of information sharing,  $\gamma$ , i.e.,  $\partial t / \partial \gamma = 0$ . Further, with respect to the effectiveness of information sharing,  $\gamma$ , a hacker's investment in cyber attacks is decreasing at a decreasing rate, their level of information sharing is increasing linearly and their pay-offs are increasing at an increasing rate. That is,  $\partial T / \partial \gamma < 0$ ,  $\partial^2 T / \partial \gamma^2 < 0$ ,  $\partial S / \partial \gamma > 0$ ,  $\partial^2 S / \partial \gamma^2 = 0$ ,  $\partial U / \partial \gamma > 0$ ,  $\partial^2 U / \partial \gamma^2 > 0$ .

(iv) with respect to the intensity of the joint information sharing,  $\lambda$ , *a* hacker's investment in cyber attacks is increasing at a decreasing rate, the level of information sharing between hackers is decreasing at an increasing rate and their pay-offs are decreasing at an increasing rate. That is,  $\partial T / \partial \lambda > 0$ ,  $\partial^2 T / \partial \lambda^2 < 0$ ,  $\partial S / \partial \lambda < 0$ ,  $\partial^2 S / \partial \lambda^2 > 0$ ,  $\partial U / \partial \lambda < 0$ ,  $\partial^2 U / \partial \lambda^2 > 0$ .

(v) with respect to the value of the firm's asset, *r*, the firms' investment in cyber security and the hacker's investment in cyber attacks are increasing linearly. That is,  $\partial t / \partial r > 0$ ,  $\partial^2 t / \partial r^2 = 0$ ,  $\partial T / \partial r > 0$ ,  $\partial^2 T / \partial r^2 = 0$ . Moreover, with respect to the value of the firm's asset, while the hackers' pay-offs increase linearly, their level of information sharing does not change. That is,  $\partial S / \partial r = 0$ ,  $\partial U / \partial r > 0$ ,  $\partial^2 U / \partial r^2 = 0$ .

(vi) the fraction of the firm's asset accruing to each hacker increases at a decreasing rate in the level of the firm's inefficiency in cyber defense and decreases at an increasing rate with its own inefficiency. That is,  $\partial K / \partial c > 0$ ,  $\partial^2 K / \partial c^2 < 0$ ,  $\partial K / \partial \theta < 0$ ,  $\partial^2 K / \partial \theta^2 > 0$ . Conversely, the firm's fraction decreases at an increasing rate in its own inefficiency while it increases at a decreasing rate in the hackers' inefficiency. That is,  $\partial k / \partial c < 0$ ,  $\partial^2 k / \partial c^2 > 0$ ,  $\partial k / \partial \theta > 0$ ,  $\partial^2 k / \partial \theta^2 < 0$ .

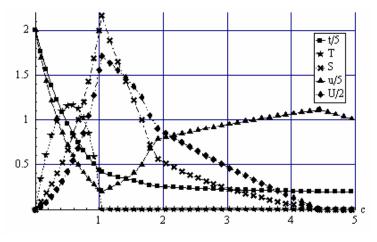


Figure 1: The variables t, T, S, u, U as functions of the firm's unit cost c.

Figs. 1-5 illustrate Propositions 1-3 while also accounting for the corner solutions, with parameter values being set such that  $c = \gamma = \lambda = 0.5$ ,  $\theta = 1$ , r = 10,  $\phi_1 = \phi_2 + \phi_3$ , except in that figure where that parameter varies. Division with 5 and 2 are done for scaling purposes. As the firm's unit cost of security investment *c* (or the level of inefficiency) increases from zero in Figure 1, the firm's investment in cyber security decreases convexly, while the hacker's investment in attack becomes inverse U shaped. A very low *c* implies that the firm's security investment is highly efficient which reduces any incentives for the hacker to attack. On the other hand, a high value of *c* causes the hacker to cut back on its attack level since the firm's defense is so modest.

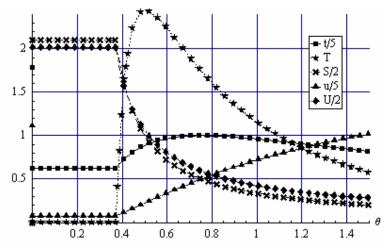


Figure 2: The variables t, T, S, u, U as functions of the hackers' unit cost  $\theta$ .

Figure 2 shows how the variables change as a function of the unit cost of attack  $\theta$ . The level of attack investment of the hackers, *T*, is low when  $\theta$  is high, and is inverse U shaped otherwise. The firm's defense level t also behaves in a similar manner. As  $\theta$  decreases, information sharing increases, and so does the hackers' profit levels. On the other hand, the firm's profit decreases as  $\theta$  decreases due to the intensity of the attack. As  $\theta$  decreases below a critical value given by  $\theta = 0.37$ , the attack level *T* decreases to zero and information sharing takes over. Note that since  $\theta$  is present only in the first order condition for *T* in (7), and not present in (6) and (8), the four variables *t*, *S*, *u*, *U* are constant when  $\theta < 0.37$ .

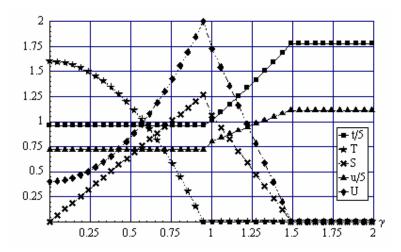


Figure 3: t, T, S, u, U as functions of the hackers' information sharing effectiveness parameter  $\gamma$ .

Figure 3 depicts the dependence on  $\gamma$  which measures how effective information sharing is relative to the attack in the composition of the overall attack. When  $\gamma = 0$ , there is no information sharing since it is not beneficial, but the attack is substantial. As  $\gamma$  increases, information sharing increases linearly and the attack level decreases convexly, eventually reaching zero as in Figure 1. The firm's defense level and profit is independent of  $\gamma$  as specified in Proposition 2. As  $\gamma$  increases above  $\gamma = 0.95$ , where T = 0, the development is similar to Figure 1, where S decreases to zero, after which t and u are constant.

Also, we find that information sharing increases in a convex manner in c, and compensates for the decrease in attack levels as c increases. As T reaches zero in accordance with Proposition 3, we find that for c=1.04, a new phenomenon emerges which is determined by solving the third order equation set. The hackers start to free ride on each other's information sharing so that the optimal level of information sharing S decreases. Also, the optimal level of investment by a firm in its security investment t decreases in the unit cost, but more moderately, and the firm's profit increases until S=0. Interestingly, Gordon et al. (2003) who focus on the cost side effects of information sharing on the overall level of security, also demonstrate similar free riding effects in information sharing using a very different model. The free-rider dilemma with respect to security investments is further analyzed by Anderson (2001).

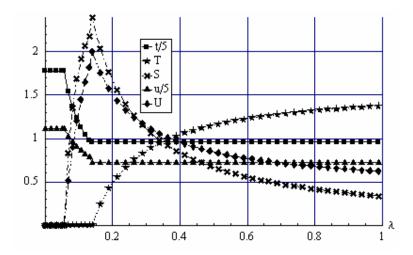


Figure 4: The variables t, T, S, u, U as functions of the intensity of joint information sharing  $\lambda$ .

Figure 4 delineates the dependence on the intensity  $\lambda$  of joint information sharing. Substantial free riding in information sharing emerges as  $\lambda$  increases, causing *S* to decrease in a concave manner. The hackers' compensate by increasing the attack *T* concavely, preserving the overall attack level  $T_A$  constant as in Proposition 1. This also ensures that the firm's security investment and pay off remains constant. The free riding among the hackers causes their profit U to decrease concavely. As  $\lambda$  decreases below  $\lambda = 0.14$  where T = 0, information sharing and hacker profit decrease to zero when  $\lambda = 0.06$ , below which *t* and *u* are constant.

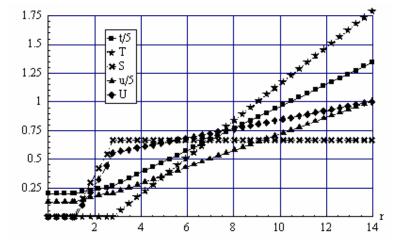


Figure 5: The variables *t*, *T*, *S*, *u*, *U* as functions of the firm's asset *r*.

Figure 5 plots the dependence on the firm's asset r. Note that all the variables increase linearly in r except information sharing, S, which is independent of r. That is, a firm's cyber security investment level, the hackers' cyber attack levels and their pay-offs increase in the value of the

firm's asset. As the asset value decreases below a critical value given by r = 2.78, the hackers do not find investing in an attack worth while, and consequently set their attack levels T = 0. As r decreases below a critical level given by r = 1.12, information sharing levels, S, and hacker profit, U, decrease to zero, while the firm's investment in cyber security, t, and its pay-off, u, are constant, similar to Figure 4. In sum, Figures 2, 4 and 5 provide some common insights. Similarly, Figures 1 and 3 provide some common insights

## 4. Endogenizing the firm's asset

Considering the asset r as exogenously given is common in the rent seeking literature. However, rents and assets frequently have to be produced. Furthermore, since firms have limited resources, they make a tradeoff between producing the asset and investing in cyber security to defend it. This calls attention for the need to examine cases where the firm decided to endogenize the asset.

Hausken (2005) has compared and contrasted rent seeking models and production and conflict models. To consider the latter assume that the firm has a resource R (e.g. a capital good, or labor) which can be divided into effort p to produce (build) the asset, and security investment t to defend or protect the asset from attack. With unit conversion costs a and b, respectively, of transforming the resource into defending versus producing the asset, the budget constraint for the firm is

$$R = at + bp \tag{11}$$

Assume that the production function for the asset takes the following simple form,

$$p = (R - at)/b \tag{12}$$

which follows from solving (11) with respect to p, and which means a linear production function. Substituting r with p in (4), and removing ct since the defense expenditure of security investment is now endogenized, (4) becomes

$$u = \frac{t}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j} (R - at) / b - \left[ (\phi_1 - \phi_2)(S_i^2 + S_j^2) - 2\phi_3 S_i S_j \right],$$

$$U_i = \frac{T_i + \gamma S_j + (\lambda/2)S_i S_j}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j} (R - at) / b - \theta T_i,$$

$$U_j = \frac{T_j + \gamma S_i + (\lambda/2)S_i S_j}{t + T_i + \gamma S_j + T_j + \gamma S_i + \lambda S_i S_j} (R - at) / b - \theta T_j$$
(13)

Calculating  $\partial u / \partial t = 0$ ,  $\partial U_i / \partial T_i = 0$ ,  $\partial U_i / \partial S_i = 0$ , and thereafter setting  $T = T_i = T_j$  and  $S = S_i = S_j$  in equilibrium, gives the three first order conditions given as follows:

FOC 
$$t: 2R(T+\gamma S) + R\lambda S^2 - at(t+4T+2S(2\gamma+\lambda S)) = 0$$
 (14)

FOC 
$$T: \frac{(R-at)(2(t+T+\gamma S)+\lambda S^2)}{2b(t+2(T+\gamma S)+\lambda S^2)^2} - \theta = 0$$
 (15)

FOC 
$$S: 2\gamma(T+\gamma S) + \lambda S(\gamma S - t) = 0$$
 (16)

Solving the above first order conditions with respect to *t*, *T*, *S*, gives us the following interior solution:

$$t = \frac{4R}{7a + 2b\theta + \sqrt{a^2 + 12ab\theta + 4b^2\theta^2}} > 0,$$

$$T = \left[-a\gamma^2(a^3 + 15a^2b\theta + 48ab^2\theta^2 + 12b^3\theta^3) + bR\theta(a^2 + 8ab\theta - 4b^2\theta^2)\lambda + \sqrt{a^2 + 12ab\theta + 4b^2\theta^2}(-a\gamma^2(a^2 + 9ab\theta + 10b^2\theta^2) + bR\theta(a + 2b\theta)\lambda)\right] /$$

$$\left[8b^2\theta^2(a^2 + 7ab\theta + 2b^2\theta^2 + (a + b\theta)\sqrt{a^2 + 12ab\theta + 4b^2\theta^2}})\lambda\right] > 0,$$

$$S = \frac{\gamma\left(a - 2b\theta + \sqrt{a^2 + 12ab\theta + 4b^2\theta^2}\right)}{4b\theta\lambda} > 0$$

Confining attention to the interior solution, Figs. 6-11 illustrate how these variables change with respect to a given parameter. The other parameter values are set to similar levels as in section 3, that is  $\gamma = \lambda = 0.5$ ,  $\theta = 1$ ,  $\phi_1 = \phi_2 + \phi_3$ . Additionally consider a = b = 1 and R = 10.

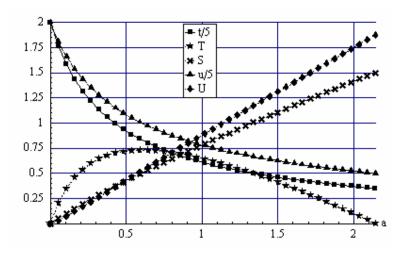


Figure 6: t, T, S, u, U as functions of the hackers' unit defense cost a.

Figure 6 shows how these variables change with respect to the unit defense cost a, and is similar to Figure 1 for the unit cost c, partly preserving the curvature. The firm's cyber security investment level, t, and its pay-off function, u, decreases convexly. On the other hand, the hackers' attack levels, T, is inverse U shaped, while their information sharing level S and pay-off functions U increase linearly.

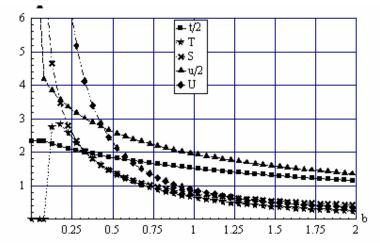


Figure 7: t, T, S, u, U as functions of the firm's unit production cost b.

Figure 7 shows the dependence on the firm's unit production cost b. An increase in b is uniformly detrimental for all the variables causing them to decrease in a convex manner with the exception of T which is inverse U shaped. The reason is that when b is very low, the firm can easily increase it security investment t because such investments are highly efficient. This causes the hackers to reduce their attack investment levels and rely on higher levels of information sharing instead.

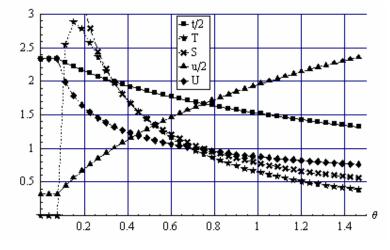


Figure 8: t, T, S, u, U as functions of the hackers' unit cost  $\theta$ .

Figure 8 shows the dependence on the unit cost of attack  $\theta$ , and is similar to Figure 2 for 4 of the 5 variables, preserving the curvature. That is, *T* is inverse U shaped, *u* increases concavely, and *S* and *U* decrease convexly. The one exception is t which decreases in Figure 8 and is inverse U shaped in Figure 2. Differences like that are expected when endogenizing the firm's asset.

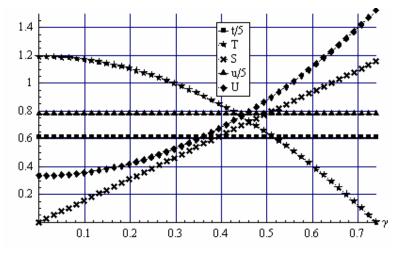


Figure 9: t, T, S, u, U as functions of the information sharing effectiveness y.

Figure 9 shows the dependence on  $\gamma$  which measures the effectiveness of information sharing relative to the attack in the composition of the overall attack. Figure 9 is similar to Figure 3, and preserve a similar curvature for the different variables. That is, *t* and *u* are constant, while *T* decreases concavely, *S* increases linearly, and *U* increases convexly.

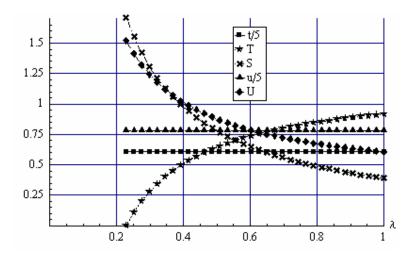


Figure 10: *t*, *T*, *S*, *u*, *U* as functions of the intensity of joint information sharing  $\lambda$ .

Figure 10 shows the dependence on the intensity  $\lambda$  of joint information sharing. Figure 10 is similar to Figure 4, preserving the curvature. That is, *t* and *u* are constant, *T* increases convexly, and *S* and *U* decreases convexly.

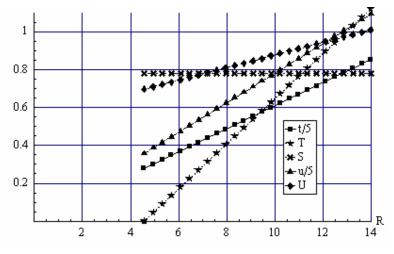


Figure 11: *t*, *T*, *S*, *u*, *U* as functions of the firm's resource *R*.

Figure 11 shows the dependence on the firm's resource R. Note that similar to Figure 5, all the variables increase linearly in r except information sharing, S, which is independent of r.

Summing up, we find that endogenizing the firm's asset largely confirms the analysis with a fixed asset in Section 3. An increase in the unit production cost of the firm is detrimental for all the five variables except the hackers' attack level which is inverse U shaped.

#### 5. An upper bound on each hacker's fraction of the asset

In this section we discuss an extension where the hackers may not want to exceed a certain limit with regard to the proportion of the firm's asset that they acquire. This can happen when hackers are apprehensive of undue attention and visibility that can cause them to be proactively targeted by the firm. If the hacker is subject to enforcement, it would not enjoy any benefit from targeting the firm's asset beyond a certain limit. One of the key phenomena that has emerged over the last five years is that various kinds of software tools or code used for hacking (such as malware embedded in consumer software) are now being written for explicit financial rewards rather than for other purposes such as reputation. Indeed, a growing trend is that of a parallel economy of hackers, quite in the spirit of Adam Smith's division of labor, in which some hackers specialize in tool creation, others trade vulnerabilities, others buy and sell credit card numbers, and then of course there are those who engage in spamming and phishing. This is starting to assume many of the aspects of the above-ground market, including externalities. For example, pay-TV operators now wait until one bad agent has established itself as the monopoly provider of forged smartcards (first-to-market usually removes incentives for further investment in forgery as it can undercut later competitors). Once this agent has x % of the market, which may be 5% or some other percentage dependent on the pay-TV operators tolerance and benefitcost analysis, they will close it down or the government can shut it down.<sup>9</sup>

Just as it is virtually impossible to secure a computer against the most persistent hacker, it is also virtually impossible for a serial hacker to avoid detection and capture. No one can expect to succeed against all opponents forever. Indeed it has been documented that "cyber detectives" are out there laying traps for and ultimately apprehending "cyber criminals"(Kremen 1998). Besides the US, a large number of countries have adopted statutes designed to protect electronic commerce and information stored on computers internally. In many instances these countries cooperate with the United States in order to apprehend hackers operating inside their borders (Platt 1996). These laws can increases the penalties from being caught and as such may reduce the hackers' incentives to grab more than a certain fraction of the target's assets. This section provides some preliminary insights into the profits of the hacker and the firm in such situations.

<sup>&</sup>lt;sup>9</sup> We thank Ross Anderson for pointing this out and suggesting this example.

In our model, recall from (9) that the fraction of the firm's asset accruing to each hacker is given by  $K = c/(2\theta+c)$  when both hackers are symmetric in their capabilities. Let Q denote the threshold value of the firm's asset acquired by each hacker beyond which the hackers are actively traced and targeted by enforcement authorities in order to be closed down. This could be an external agent like a government institution, or it could be a division within the firm itself which operates under other incentives and constraints than those assumed in the model in this paper. Thus, if the fraction of the firm's asset accruing to each hacker is larger than this constraint, i.e., if K>Q, where the firm takes into account how many hackers operate near this threshold, each hacker exceeding the threshold gets closed down in some manner.

Assume that the threshold Q, and also the inefficiencies c and  $\theta$  of defense and attack, are common knowledge. The new optimization problem for each hacker then becomes as follows. If K < Q, which means that the firm's threshold for being hacked is so large that it accepts that a fraction K of its asset r is confiscated by each hacker, then the solution in (9) and the discussion in the previous sections apply. Conversely, if K > Q, each hacker gets closed down if it is not intelligent, earning zero profit as its investment and information sharing eventually decrease to zero. (At the moment when it gets closed down, no longer earning a fraction of the asset, its profit is negative determined by its cost expenditure  $\theta T$  where T is given in (9).) If K > Q and each hacker is intelligent, it chooses to hack a fraction  $K = Q - \varepsilon$ , where  $\varepsilon > 0$  is arbitrarily small but positive, which means that it hacks marginally less than the threshold level which causes its elimination. Since there are two hackers, these together earn the fraction  $2(Q - \varepsilon)$ , while the firm retains the remaining fraction  $1-2(Q - \varepsilon)$ . Using (3), (4) and (10), the fractions earned by the firm and each hacker are

$$k = \frac{t}{t + 2T_A} = 1 - 2(Q - \varepsilon), \quad K = \frac{T_A}{t + 2T_A} = \frac{1 - k}{2} = Q - \varepsilon$$
(18)

This implies that the cumulative attack level is given by

$$T_{A} = \frac{Q - \varepsilon}{1 - 2(Q - \varepsilon)}t\tag{19}$$

Suppose there are no leakage costs, i.e.,  $\phi_1 = \phi_2 + \phi_3$ . Consider first the case when the firm decides to keeps its investment level *t* the same as in (9). Then, the firm's profit is given by

$$u = kr - ct = \left(1 - 2(Q - \varepsilon) - \frac{2c(2\theta - c)}{(2\theta + c)^2}\right)r$$
(20)

If each hacker chooses its overall attack  $T_A$  to consist of security investment only, and no information sharing, that is,  $T_A = T$  and S=0, its profit is given by

$$U = Kr - \theta T = \left(1 - \frac{2\theta(2\theta - c)}{(2\theta + c)^2 [1 - 2(Q - \varepsilon)]}\right) (Q - \varepsilon)r$$
(21)

Two main changes are possible to this solution where an intelligent hacker hacks a limited fraction  $K = Q - \varepsilon$  to avoid being caught. First, the firm may choose an alternative level of investment *t* than the solution in (9). That is, the firm may realize that each of two hackers are intelligent and will choose  $T_A = T$  as determined in (19) which is dependent on *t*. Since the fractions accruing to each hacker are fixed and predetermined as in (18), this becomes a coordination game where the firm and each hacker adjust their investments *t* and *T* so that (18) is satisfied. To minimize their costs, one solution is that both parties reduce their investments *t* and *T* toward zero while (18) is satisfied, though that solution seems hard to occur in practice.

The other possible change to the solution in (20) and (21) is that the hacker chooses another allocation between security investment and information sharing instead of setting  $T_A = T$  and S=0. A simple solution is to assume that the hackers agree that investment is costly while information sharing is not, and hence choose only to share information. That is,  $T_i=T_j=T=0$ . Setting the left hand side of (10) equal to the right hand side of (19) gives

$$\gamma S + (\lambda/2)S^2 = \frac{Q - \varepsilon}{1 - 2(Q - \varepsilon)}t$$
(22)

Substituting the value of  $t = \frac{2r(2\theta - c)}{(2\theta + c)^2}$  gives us the following equation

$$S = \frac{-\gamma + \frac{\sqrt{\gamma^2 (-1 + 2(Q - \varepsilon))(c + 2\theta)^2 + 4r(Q - \varepsilon)(c - 2\theta)\lambda}}{\sqrt{-1 + 2(Q - \varepsilon)}(c + 2\theta)}}{\lambda}$$
(23)

The firm's profit u remains as in (19) while each hacker's profit is given by<sup>10</sup>

$$U = Kr = (Q - \varepsilon)r.$$
<sup>(24)</sup>

<sup>&</sup>lt;sup>10</sup> A possible alternative to the method in this section may be to perform some kind of a constrained optimization with Lagrangian multipliers.

#### 6. Extensions and limitations

In our analysis, we have only considered symmetric equilibrium, implying that both hackers have equal characteristics and thus are equally efficient in utilizing the shared security information in attacking firms. A future extension could involve looking at asymmetric situations between hackers, where hackers have unequal characteristics in various respects. Asymmetry is analytically more challenging, with less likelihood of analytical solutions for the most general scenarios. Our results are robust to some extensions such as changes in timelines of the game. For example, it is possible to envisage a scenario in which all three actors first chose security attacks  $T_{i}$ ,  $T_{j}$  and t, and then in the next stage, the two hackers choose sharing levels  $S_{i}$ ,  $S_{j}$  simultaneously. The notion being that cyber security investments are long term decisions, and hence are chosen before information sharing decisions which are in principle easy to change. Our preliminary analysis reveals that there is no major qualitative change in results in such a game.

Further, another limitation is that we have only considered a scenario with two hackers. Generalizing this research to include *n* hackers could be an interesting extension but becomes analytically tedious in our setup. However, we conjecture that many of the qualitative insights of our model carry through to a situation with more than two hackers. A difference is that a firm facing more than two hackers is subject to an opposition which may share information in more sophisticated manners than what is accounted for in the current model which distinguishes between effectiveness of information sharing and intensity of joint information sharing.

Other extensions are to different kinds of security investment, and distinguishing between different kinds of information that hackers can share. Information is multidimensional. It may pertain to high level or low level security breaches, methods and success of earlier attacks, identities of hackers, and secrets about research, development, future plans, trade, capacities, personnel dispositions, etc.

Despite these limitations, we believe that our model addresses an important issue, and hope that the proposed approach may be used as a starting point for additional research in this area.

#### 7. Conclusion

The paper considers a scenario where one firm is subject to a cyber attack by two external hackers. The hackers choose the optimal level of information to share with each other about the firm's vulnerabilities and security breaches, and choose the optimal attack levels. The firm chooses the optimal defense, which is costly and consists in investing in information technology security to protect its asset. The hackers collect information in various manners, and attempt to gain access to the information the firms collect about their security breaches. Each hacker prefers to receive information from the other hacker about the firm's vulnerabilities. Providing information to the other hacker places the other hacker in a better position when competing for the firm's asset. The hackers benefit from joint information sharing. The paper analyzes the extent to which a hacker has incentives to provide information and investing in an attack, which is costly. Each hacker thus has two free choice variables, and the firm has one free choice variable.

The paper shows that each hacker's attack and information sharing are strategic complements such that increasing one causes an increase in the other, and vice versa. Conversely, one hacker's attack and the other hacker's information sharing are strategic substitutes such that increasing one causes a decrease in the other, and vice versa.

The paper defines an overall attack for the hacker which is a weighted sum of three components: (i) the investment in an attack, (ii) the amount of information sharing, and (iii) the amount of joint information sharing. Each of the two hackers substitutes between these three components such that the relative weights of the second and third components do not, remarkably, affect the defense security investment of the defending firm, and also do not affect the fractions of the asset accruing to each hacker and the firm. For the special case of no leakage costs, the firm's profit is also unaffected by the relative weights of the second and third components. The firm is affected by the hackers' overall attack level, and not their tradeoffs between the three components. The overall investment in attack increases concavely in the firm's unit defense cost, decreases convexly in the hackers' unit cost of attack, and increases linearly in the firm's asset.

As the firm's unit cost of security investment increases from zero towards twice the hackers' unit cost of attack, the attack follows an inverse U shaped function and reaches zero while the firm's security investment and the hackers' information sharing levels remain positive. Specifically, while the firm's cyber security investment and profit decrease convexly, the hackers' information sharing increases convexly, and the hackers' profits increase as well.

As the hackers' unit cost of attack increases, the firm's profit increases, while the hackers' information sharing and profit decrease convexly. The firm's defense and the hackers' attack are inverse U shaped for some regions in the parameter space. As the effectiveness of information sharing relative to the investment in cyber attack in the composition of the overall attack increases, the firm's defense and profit are constant, the hacker's attack levels decrease convexly, their information sharing increases linearly, and their profit increases convexly.

As the intensity of joint information sharing increases, the firm's defense and profit are constant, the attack increases concavely, and the hackers' information sharing and profit decrease convexly. Increasing the firm's asset causes all the variables to increase linearly, except information sharing which is constant. We also find that endogenizing the firm's asset largely confirms the analysis with a fixed asset. Increasing the unit production cost is detrimental for all the five variables except the hackers' attack which is inverse U shaped.

A few policy implications and managerial insights of the analyses merit some discussion. First, our analysis also reveals that the cumulative attack level of the hackers is not affected by the effectiveness of information sharing between them and moreover, is also unaffected by the intensity of joint information sharing. If it turns out that information sharing between hackers does not enhance their attack intensities, then there is possibly less cause for concern from such information sharing than it is thought to be. This result can have some ramifications on firms' incentives and strategies regarding their adversarial contest with hackers especially in situations involving industrial espionage which is widely being touted as new mantra for hackers. It also undermines to some extent, the importance of the truthful revelation of information by hackers to each other.

Second, our analysis suggests the need to heighten firm's awareness that hackers not only choose strategically how much to invest in an attack, and that hackers may compete with each

other in attacking more successfully, but also that hackers may cooperate through sharing information with each other about a firm's vulnerabilities. Firms may defend themselves by designing their security investments such that when it is breached by one hacker, the sensitive information may quickly become obsolete which limits the potential vulnerability if the information is transferred elsewhere. Also, if a firm knows the identity of the hacker that has breached its security, it may focus on not only blocking future attacks by this hacker, but also on providing incentives to the hacker so that the sensitive information does not get transferred to other hackers. Such incentives can be of monetary, informational, political, persuasional, or disinformational nature.

Finally, if a hacker's information sharing can be reduced, strategic complementarity with the hacker's attack also reduces the latter. This suggests compartmentalization, in some sense or other, the interaction of hackers with each other or designing incentives to isolate them from each other rather than treating them as one monolithic group which acts in unison. Moreover, the fact that the intensity of information sharing and effectiveness of sharing have opposite impact on hackers' investment in attacks, their information sharing levels and profits provides further insights into mechanisms that firms can design to take preventive measures and thwart cyber attacks.

### References

Anderson, R., 2001. Why information security is hard: An economic perspective," *Proceedings* of 17th Annual Computer Security Applications Conference, December.

Arora, A., R. Krishnan, R. Telang, Yang, Y., 2005. An empirical analysis of vendor response to software vulnerability disclosure. Working Paper, Carnegie Mellon University, August 2005.

Cavusoglu, H., B. Mishra, B., Raghunathan, S., 2005. The value of intrusion detection systems in information technology security architecture. *Information Systems Research* 16, 1, 28-46.

Choi, J., C. Fershtman, Gandal, N., 2005. The economics of internet security. Department of Economics, Michigan State University, December 6, 2005.

Gal-Or, E., 1985. Information sharing in oligopoly. *Econometrica* 53 (2), 329–343.

Gal-Or, E., Ghose, A., 2003. The economic consequences of sharing security information. In: *Proceedings of the Second Workshop on Economics and Information Security*, May 29-30, University of Maryland.

Gal-Or, E., Ghose, A., 2005. The economic incentives for sharing security information. *Information Systems Research* 16 (2), 186-208.

Gordon, L.A., Loeb, M., 2001. Using information security as a response to competitor analysis systems. *Communications of the ACM* 44, 9, 70-75.

Gordon, L.A., Loeb, M., 2003. Expenditures on competitor analysis and information security: A managerial accounting perspective. In Bhimani, A. (ed.), Management Accounting in the New Economy, Oxford University Press, 95-111.

Gordon, L.A., Loeb, M., Lucyshyn, W., 2003. Sharing information on computer systems security: An economic analysis. *Journal of Accounting and Public Policy* 22 (7), 461-485.

Gordon, L.A., Loeb, M., Lucyshyn, W., Richardson, R., 2004. 2004 CSI/FBI computer crime and security survey. *Computer Security Journal* XX (3), 33-51.

Hausken, K., 2005. Production and conflict models versus rent seeking models, *Public Choice* 123, 59-93.

Hausken, K., 2006. Information sharing among firms and cyber attacks. Paper presented May 24, 2006 at the Annual Forum on "Financial Information Systems and Cyber Security: A Public Policy Perspective", Robert H. Smith School of Business, University of Maryland.

Howard, J., 1997. Analysis of security incidents on the Internet. Unpublished Doctoral Dissertation, Carnegie Mellon University, www.cert.org/research/JHThesis/Start.htm .

Hirshleifer, J., 1995. Anarchy and its breakdown. Journal of Political Economy 103(1), 26-52.

Kermack, W.O., McKendrick, A., 1927. Contributions to the mathematical theory of epidemics. *Proceedings of the Royal Statistical Society* 115: 700-721.

Kirby, A., 1988. Trade associations as information exchange mechanisms. *RAND Journal of Economics* 29 (1), 138-146.

Kjaerland, M., 2005. A classification of computer security incidents based on reported attack data, *Journal of Investigative Psychology and Offender Profiling* 2, 105-120.

Kremen, H., 1998. Apprehending the computer hacker: The collection and use of evidence. *Computer Forensics Online.* 

Kunreuther, H., Heal, G., 2003. Interdependent security. *The Journal of Risk and Uncertainty* 26, 2/3, 231-249.

Lanchester, F.W., 1916. Aircraft in warfare: the dawn of the fourth arm, Tiptree Constable and Co. Ltd., London.

Lin. Y., 2003. The institutionalization of hacking practices. Ubiquity. Volume 4, Issue 4.

Nizovtsev, D., M. Thursby. 2005. Economic analysis of incentives to disclose software Vulnerabilities. Working Paper.

Novshek, W., Sonnenschein, H., 1982. Fulfilled expectations in cournot duopoly with information acquisition and release. *Bell Journal of Economics* 13 (1), 214-218.

Png, I., C. Tang, Wang, Q., 2006. Information security: User precautions and hacker targeting. Working Paper, National University of Singapore.

Platt, C. 1996. Anarchy Online (Net Crime/Net Sex), Harper Collins, New York.

Raymond, E., 2001. The cathedral and the bazaar: Musings on linux and open source by an accidental revolutionary. Revised edition. O'Reilly.

Ritchie, C., 2000. A look at the security of the open source development model. Technical Report, Oregon State University.

Risan, L., 2000. Hackers produce more than software, they produce hackers. Version 2.1 http://folk.uio.no/lr isan/Linux/Identity games/

Salop, S.C., Scheffman, D., 1983. Raising rivals' costs. A.E.R. Papers and Proceedings, 73, 267-271.

Shapiro, C., 1986. Exchange of cost information in oligopoly. *Review of Economic Studies* 53 (3), 433–446.

Tullock, G., 1967. The welfare costs of tariffs, monopolies, and theft. *Western Economic Journal* 5, 224-232.

Tullock, G., 1980. Efficient rent seeking. In: J. Buchanan, R. Tollison and G. Tullock, (Eds.), Towards a Theory of the Rent-Seeking Society, College Station, Texas A&M University Press, pp. 97-112.

Vives, X., 1990. Trade association disclosure rules, incentives to share information, and welfare. *RAND Journal of Economics* 21 (3), 409–430.

## Appendix

## **Proof of Proposition 1**

Proof: The proof follows from inserting (9) into (10) and considering the signs of the first and second derivatives of (10). These are

$$\frac{\partial T_A}{\partial c} = \frac{2r(2\theta - c)}{(2\theta + c)^3}, \\ \frac{\partial^2 T_A}{\partial c^2} = -\frac{4r(4\theta - c)}{(2\theta + c)^4}, \\ \frac{\partial T_A}{\partial \theta} = -\frac{8cr}{(2\theta + c)^3}, \\ \frac{\partial^2 T_A}{\partial \theta^2} = \frac{48cr}{(2\theta + c)^4}, \\ \frac{\partial T_A}{\partial r} = \frac{2c}{(2\theta + c)^2}$$
(A1)

# **Proof of Proposition 2**

Insert c=2 $\theta$ - $\epsilon$  into (9) where  $\epsilon$ >0 is arbitrarily small but positive. This causes the second term within the bracket in the expression for T to be arbitrarily large and positive, which means that the expression for T becomes negative with arbitrarily large absolute value. Since the hackers' attack cannot be negative, this gives T=0, t>0, S>0 for a sufficiently small  $\epsilon$ . QED

## **Proof of Proposition 3**

This follows from considering the signs of the first and second derivatives of (9).

(i)

$$\frac{\partial t}{\partial c} = -\frac{2r(6\theta - c)}{(2\theta + c)^3}, \frac{\partial^2 t}{\partial c^2} = \frac{4r(10\theta - c)}{(2\theta + c)^4},$$

$$\frac{\partial T}{\partial c} = \frac{4\gamma^2\theta(2\theta + c)^4 - 2r\lambda(2\theta - c)^4}{(c^2 - 4\theta^2)^3\lambda}, \frac{\partial^2 T}{\partial c^2} = -\frac{4r(4\theta - c)}{(2\theta + c)^4} - \frac{8\gamma^2\theta(4\theta + c)}{(2\theta - c)^4\lambda},$$

$$\frac{\partial S}{\partial c} = \frac{4\gamma\theta}{(2\theta - c)^2\lambda}, \frac{\partial^2 S}{\partial c^2} = \frac{8\gamma\theta}{(2\theta - c)^3\lambda},$$

$$\frac{\partial u}{\partial c} = -\frac{8r(2\theta - c)\theta}{(2\theta + c)^3}, \frac{\partial^2 u}{\partial c^2} = \frac{16r(4\theta - c)\theta}{(2\theta + c)^4}$$
(A2)

(ii)  

$$\frac{\partial t}{\partial \theta} = \frac{4r(-2\theta + 3c)}{(2\theta + c)^3}, \frac{\partial^2 t}{\partial \theta^2} = \frac{16r(2\theta - 5c)}{(2\theta + c)^4},$$

$$\frac{\partial S}{\partial \theta} = -\frac{4c\gamma}{(2\theta - c)^2\lambda}, \frac{\partial^2 S}{\partial \theta^2} = \frac{16c\gamma}{(2\theta - c)^3\lambda},$$

$$\frac{\partial U}{\partial \theta} = -c^2 \left(\frac{4r}{(2\theta + c)^3} + \frac{8\gamma^2\theta}{(2\theta - c)^3\lambda}\right), \frac{\partial^2 U}{\partial \theta^2} = 8c \left(\frac{3cr}{(2\theta + c)^4} + \frac{c\gamma^2(4\theta + c)}{(2\theta - c)^4\lambda}\right)$$
(A3)

$$\frac{\partial t}{\partial \gamma} = 0,$$

$$\frac{\partial T}{\partial \gamma} = -\frac{8c\gamma\theta}{(2\theta - c)^2\lambda}, \quad \frac{\partial^2 T}{\partial \gamma^2} = -\frac{8c\theta}{(2\theta - c)^2\lambda},$$

$$\frac{\partial S}{\partial \gamma} = \frac{2c}{(2\theta - c)\lambda}, \quad \frac{\partial^2 S}{\partial c^2} = 0,$$

$$\frac{\partial U}{\partial \gamma} = \frac{8c\gamma\theta^2}{(2\theta - c)^2\lambda}, \quad \frac{\partial^2 U}{\partial c^2} = \frac{8c\theta^2}{(2\theta - c)^2\lambda}$$
(A4)

$$\frac{\partial t}{\partial \lambda} = 0,$$

$$\frac{\partial T}{\partial \lambda} = \frac{4c\gamma^2\theta}{(2\theta - c)^2\lambda^2}, \frac{\partial^2 T}{\partial \lambda^2} = -\frac{8c\gamma^2\theta}{(2\theta - c)^2\lambda^3},$$

$$\frac{\partial S}{\partial \lambda} = -\frac{2c\gamma}{(2\theta - c)\lambda^2}, \frac{\partial^2 S}{\partial \lambda^2} = \frac{4c\gamma}{(2\theta - c)\lambda^3},$$

$$\frac{\partial U}{\partial \lambda} = -\frac{4c\gamma^2\theta^2}{(2\theta - c)^2\lambda^2}, \frac{\partial^2 U}{\partial \lambda^2} = \frac{8c\gamma^2\theta^2}{(2\theta - c)^2\lambda^3}$$
(A5)

(v)

$$\frac{\partial t}{\partial r} = \frac{2(2\theta - c)}{(2\theta + c)^2}, \frac{\partial T}{\partial r} = \frac{2c}{(2\theta + c)^2}, \frac{\partial U}{\partial r} = \frac{c^2}{(2\theta + c)^2}$$
(A6)

(vi)

$$\frac{\partial K}{\partial c} = \frac{2\theta}{(2\theta+c)^2}, \frac{\partial^2 K}{\partial c^2} = -\frac{4\theta}{(2\theta+c)^3},$$
$$\frac{\partial K}{\partial \theta} = -\frac{2c}{(2\theta+c)^2}, \frac{\partial^2 K}{\partial \theta^2} = \frac{8c}{(2\theta+c)^3},$$
$$\frac{\partial k}{\partial c} = -\frac{4\theta}{(2\theta+c)^2}, \frac{\partial^2 k}{\partial c^2} = \frac{8\theta}{(2\theta+c)^3},$$
$$\frac{\partial k}{\partial \theta} = \frac{4c}{(2\theta+c)^2}, \frac{\partial^2 k}{\partial \theta^2} = -\frac{16c}{(2\theta+c)^3}$$

(A7)